

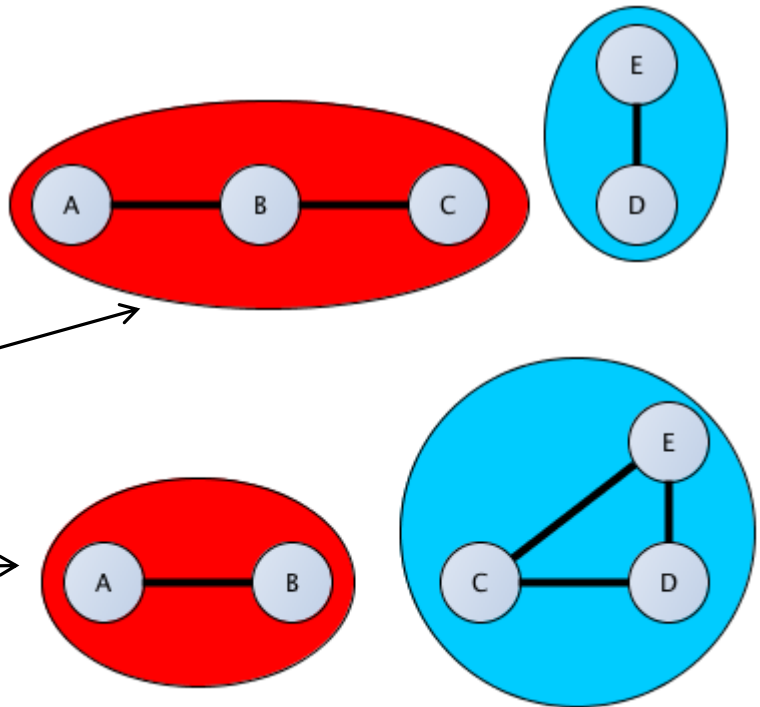
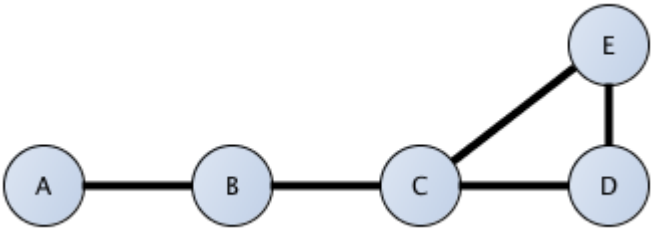
Big Data Management and Analytics Assignment 12

Assignment 12-1.1

A good partitioning fullfills two conditions:

- Maximizes the number of edges **within a group**
- Minimizes the number of edges **between groups**

Assignment 12-1.1



- Partitions shall be balanced as far as possible → the graph shall be split either in
 - $\{A, B, C\}$ and $\{D, E\}$ or
 - in $\{A, B\}$ and $\{C, D, E\}$
- The second choice has more edges within the groups, and only one edge instead of two have to be removed → Best partitioning is $\{A, B\}$ and $\{C, D, E\}$

RECAP:

The **modularity Q** of a **partitioning S** of a **graph G** is defined as follows:

$$Q \propto \sum_{s \in S} [(\#edges\ within\ group\ s) - (expected\ \#edges\ within\ group\ s)]$$

$$Q(G, S) = \frac{1}{2m} \sum_{s \in S} \sum_{i \in S} \sum_{j \in S} (a_{i,j} - \frac{k_i k_j}{2m})$$

Normalizing: $-1 < Q < 1$

Assignment 12-1.2

What we already have:

- Number of nodes $|n| = 5$
- Number of edges $|m| = 5$
- Degree of nodes:

Node	Degree
k_A	1
k_B	2
k_C	3
k_D	2
k_E	2

Assignment 12-1.2

Steps for computing Q:

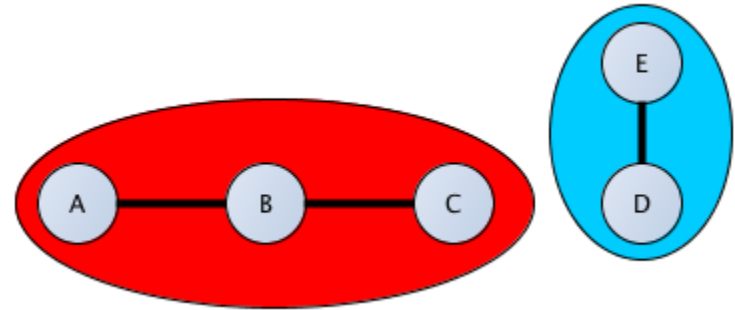
1. Compute the adjacency matrix
2. Compute the modularity matrix: $(B_{ij} = A_{ij} - \frac{k_i \cdot k_j}{2m})$
3. Sum up the entries of the single clusters
4. Sum up the sums of all clusters
5. Normalize the result

Assignment 12-1.2

Steps for computing Q:

1. Compute the adjacency matrix

$$A = \begin{pmatrix} 0 & 1 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 \end{pmatrix}$$



Assignment 12-1.2

Steps for computing Q:

2. Compute the modularity matrix: $(B_{ij} = A_{ij} - \frac{k_i \cdot k_j}{2m})$

$$B_{ij} = \frac{1}{10} \begin{pmatrix} -1 & 8 & -3 & - & - \\ 8 & -4 & 4 & - & - \\ -3 & 4 & -9 & - & - \\ - & - & - & -4 & 6 \\ - & - & - & 6 & -4 \end{pmatrix}$$

Assignment 12-1.2

Steps for computing Q:

3. Sum up the entries of the single clusters

$$s_1 = \{A, B, C\} \rightarrow \frac{1}{10}((-1) + 8 - 3 + 8 - 4 + 4 - 3 + 4 - 9) = \frac{4}{10}$$

$$s_2 = \{D, E\} \rightarrow \frac{1}{10}((-4) + 6 + 6 - 4) = \frac{4}{10}$$

$$\sum_{s \in S} = \frac{4}{10} + \frac{4}{10} = \frac{8}{10} = 0.8$$

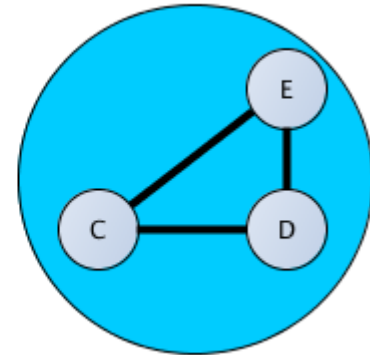
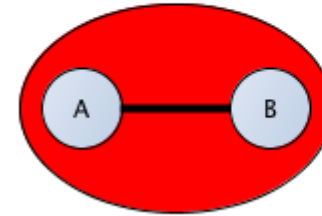
$$Q(G, S) = \frac{1}{10} \cdot (0.8) = 0.08$$

Assignment 12-1.3

Steps for computing Q:

1. Compute the adjacency matrix

$$A = \begin{pmatrix} 0 & 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 1 & 0 \end{pmatrix}$$



Assignment 12-1.3

Steps for computing Q:

2. Compute the modularity matrix: $(B_{ij} = A_{ij} - \frac{k_i \cdot k_j}{2m})$

$$B_{ij} = \frac{1}{10} \begin{pmatrix} -1 & 8 & - & - & - \\ 8 & -4 & - & - & - \\ - & - & -9 & 4 & 4 \\ - & - & 4 & -4 & 6 \\ - & - & 4 & 6 & -4 \end{pmatrix}$$

Assignment 12-1.3

Steps for computing Q:

3. Sum up the entries of the single clusters

$$s_1 = \{A, B\} \rightarrow \frac{1}{10}((-1) + 8 + 8 - 4) = \frac{11}{10}$$

$$s_2 = \{C, D, E\} \rightarrow \frac{1}{10}((-9) + 4 + 4 + 4 - 4 + 6 + 4 + 6 - 4) = \frac{11}{10}$$

$$\sum_{s \in S} = \frac{11}{10} + \frac{11}{10} = \frac{22}{10}$$

$$Q(G, S) = \frac{1}{10} \cdot \frac{22}{10} = 0.22$$

Assignment 12-1.4

- The higher the modularity Q , the better the partitioning
- Removing the edge $\{B, C\}$ yields a higher Q value than the removal of the edges $\{C, E\}$ and $\{C, D\}$
- → The hypothesis from 1 which relies on maximizing the number of edges within the groups and minimizing the number of edges between the groups was correct

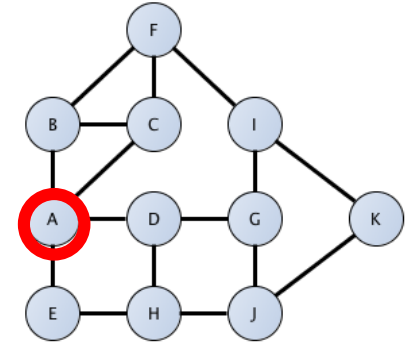
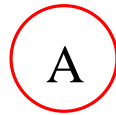
RECAP:

Girven-Newman algorithm:

1. Begin with node A and perform a BFS and construct a DAG (directed acyclic graph)
2. Count the number of shortest paths from A to all other nodes
3. Compute the **betweenness**, by traversing the tree in a bottom-up fashion. If there exist multiple paths, these are counted partially:
 1. $node\ flow = 1 + \sum childEdges$
 2. Split the flow based on the values of the parents (shortest path)

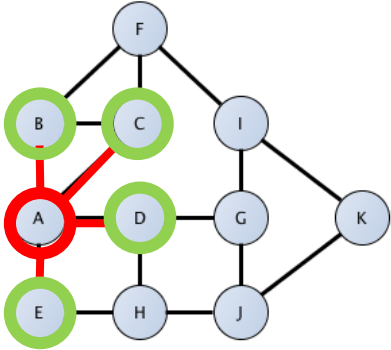
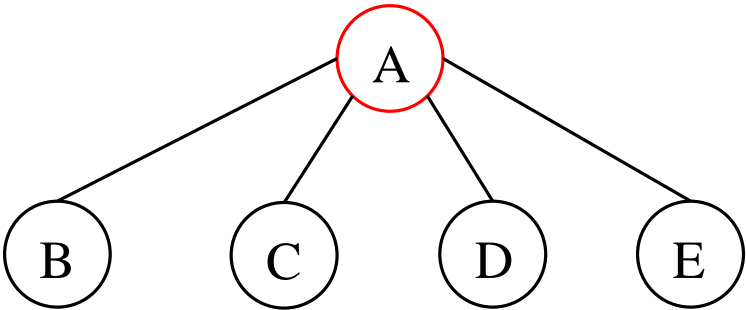
Assignment 12-2

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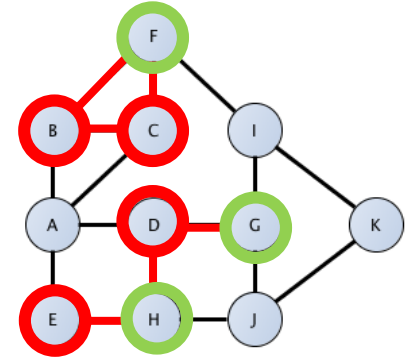
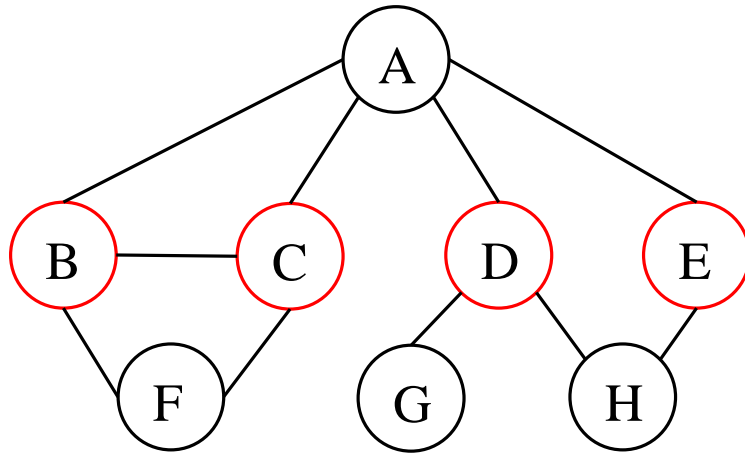
Assignment 12-2

Step 1:



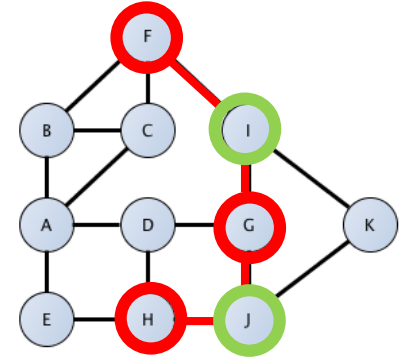
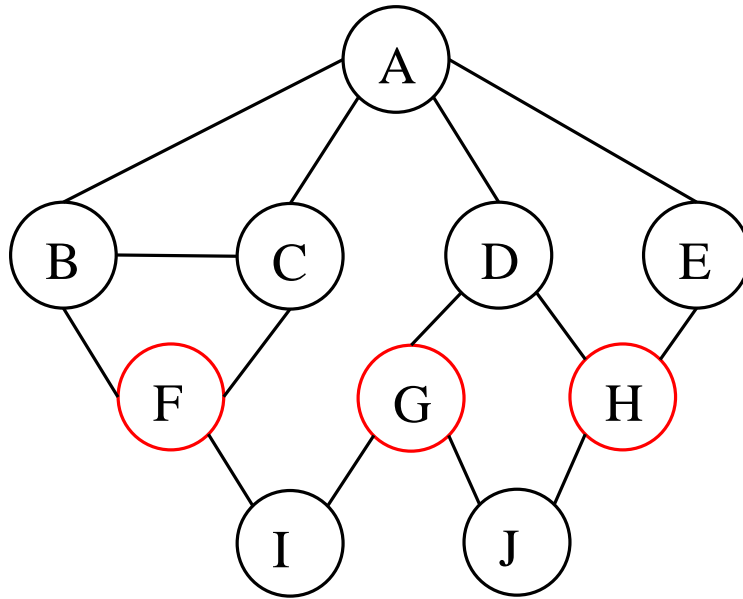
Assignment 12-2

Step 1:



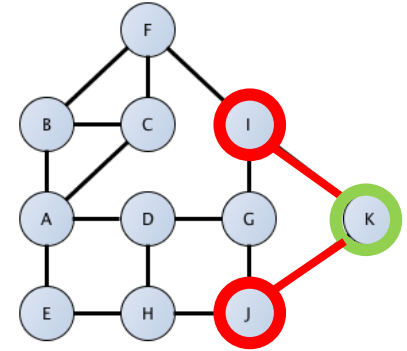
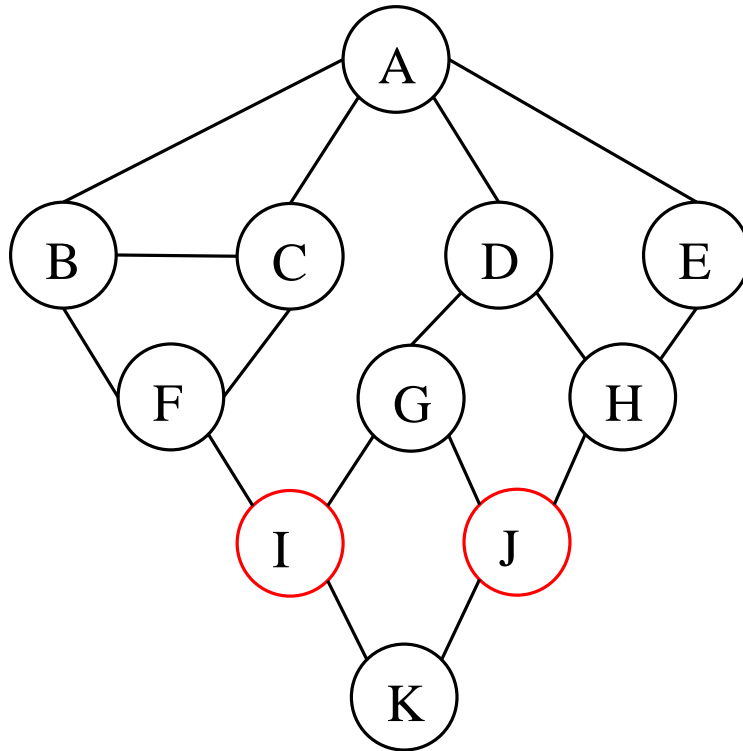
Assignment 12-2

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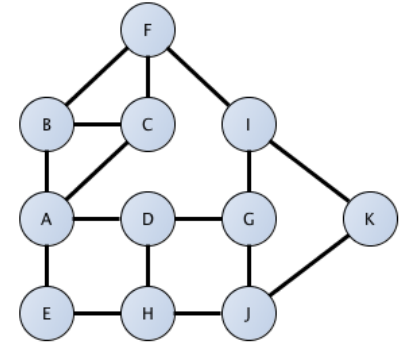
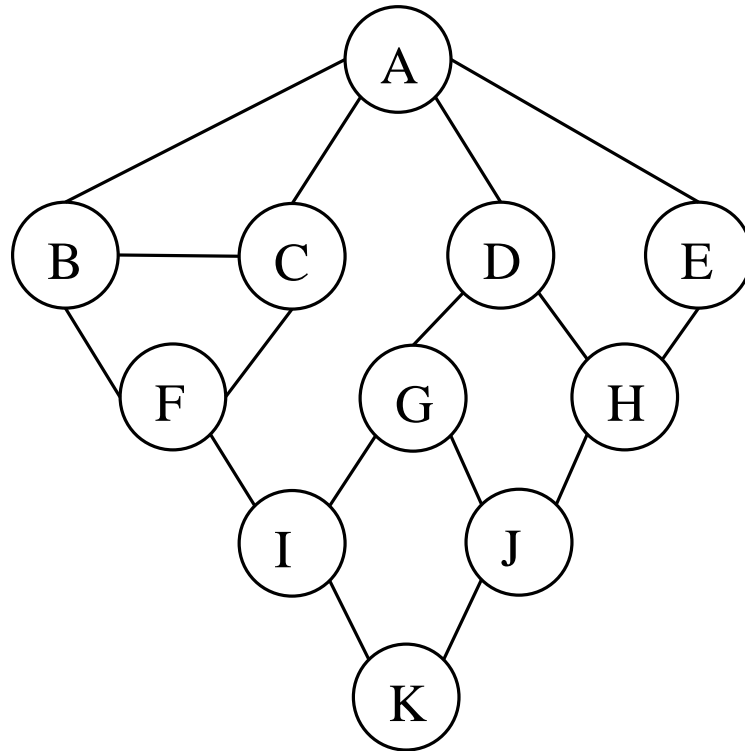
Assignment 12-2

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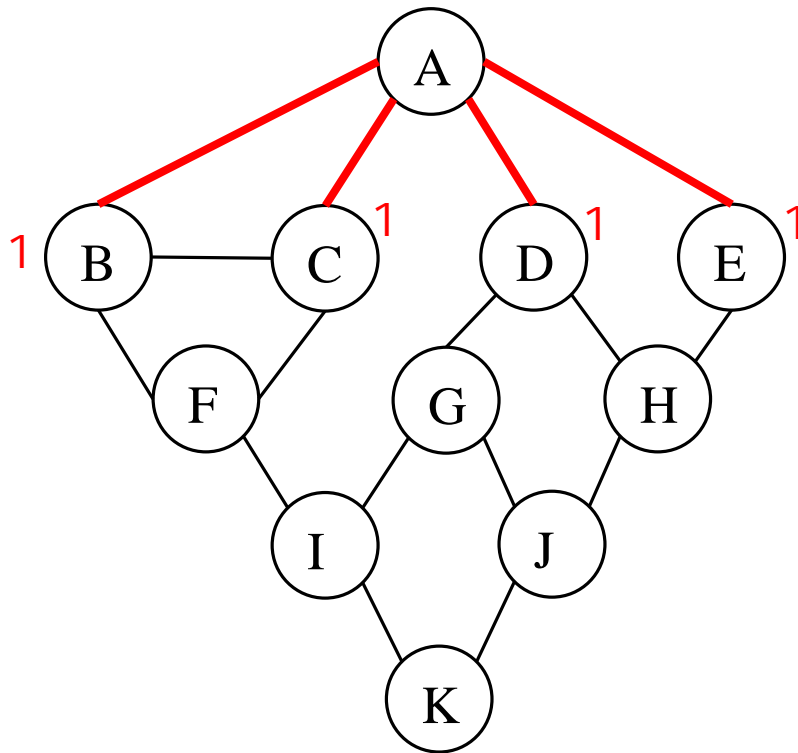
Assignment 12-2

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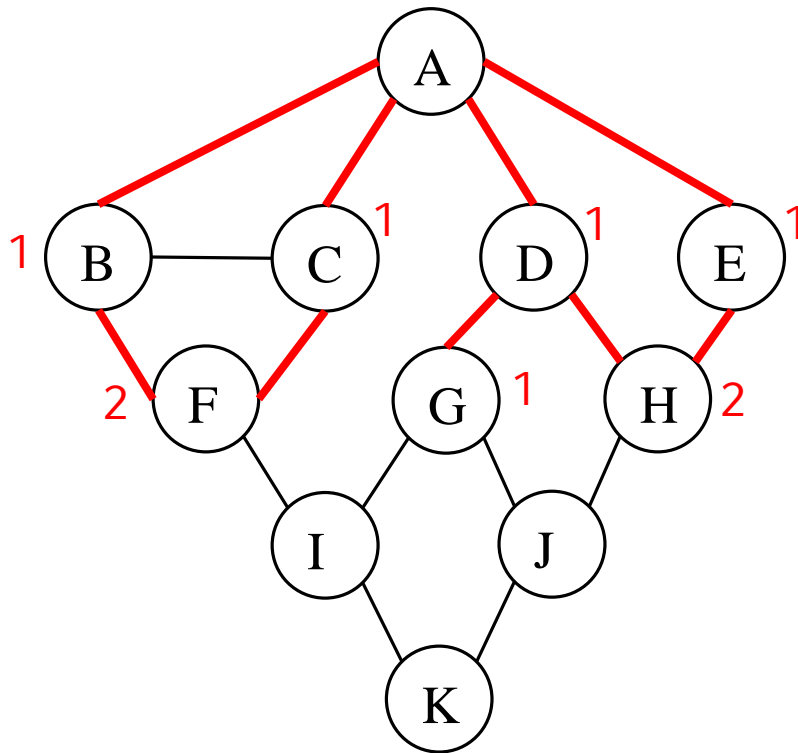
Assignment 12-2

Step 2:



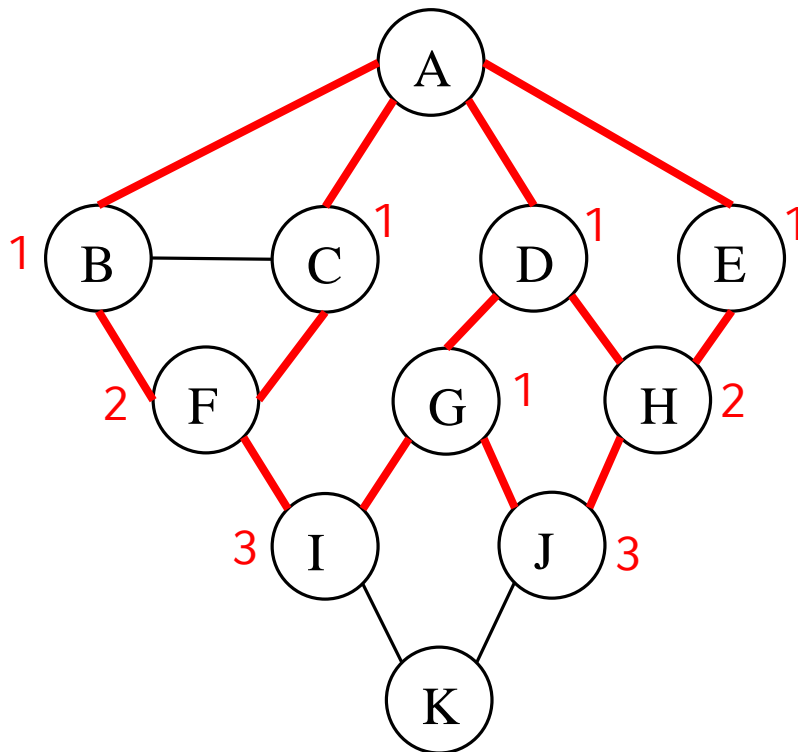
Assignment 12-2

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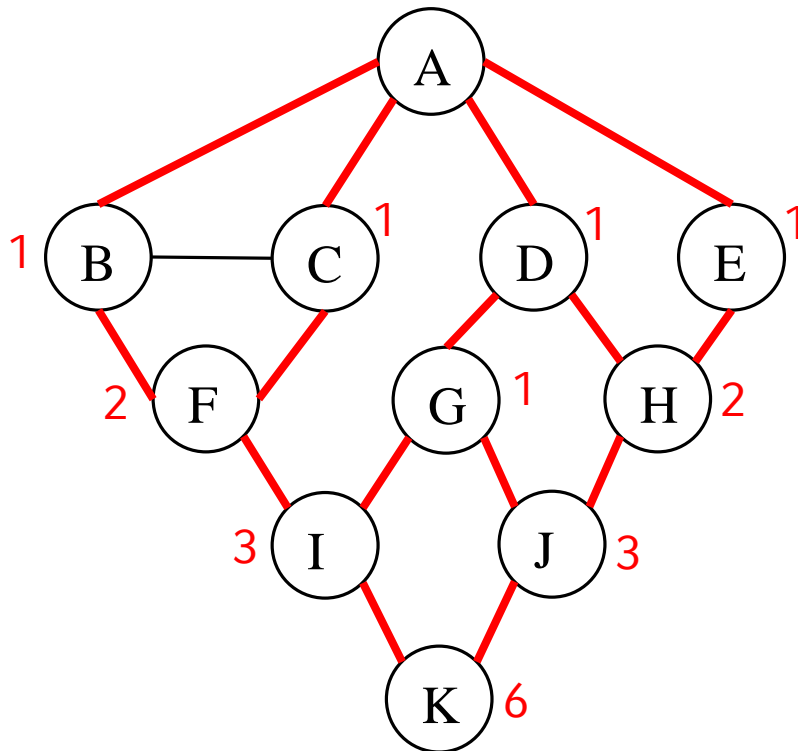
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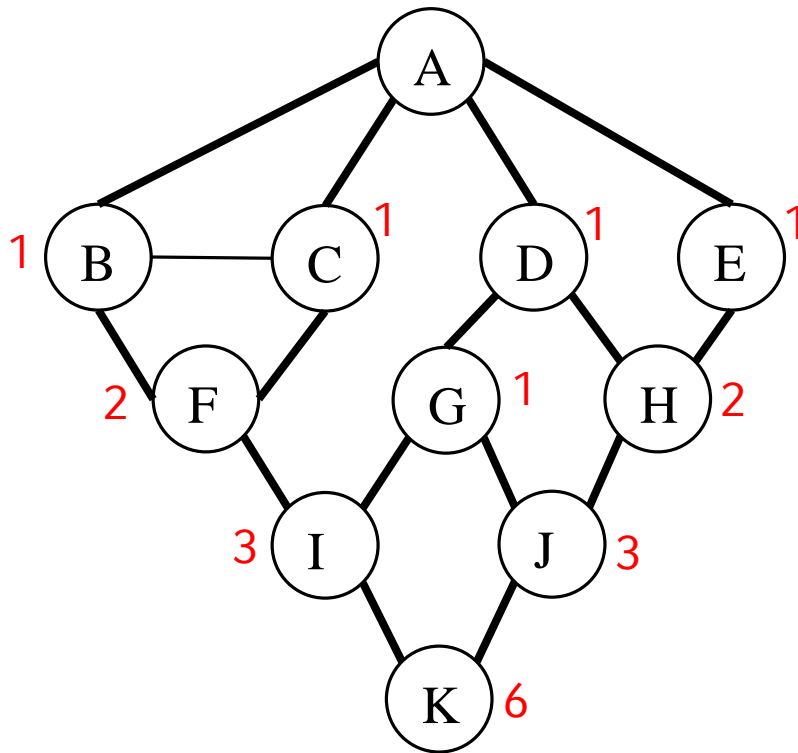
Assignment 12-2

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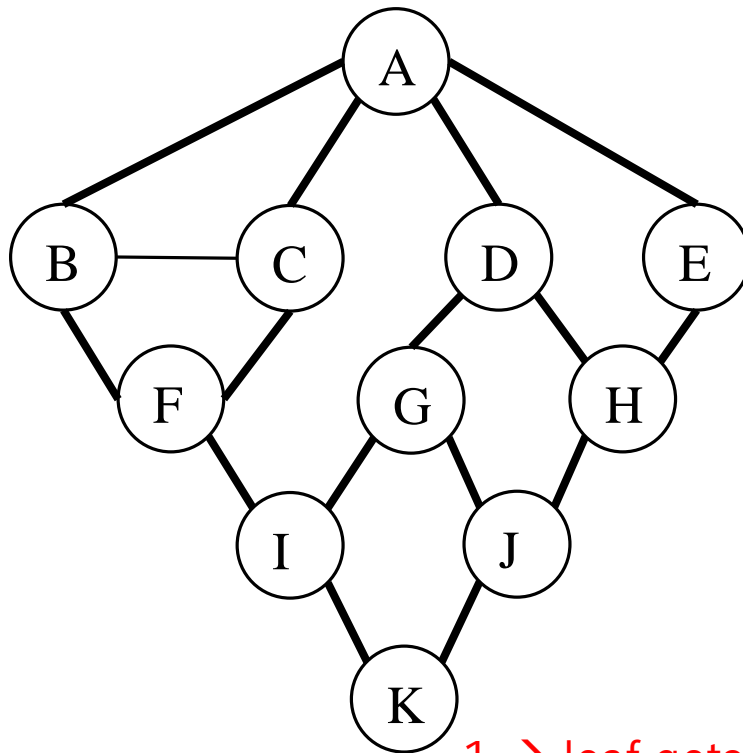
Assignment 12-2

Step 2:

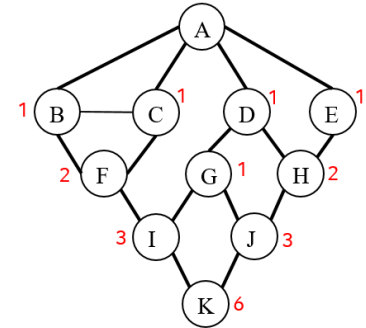


Assignment 12-2

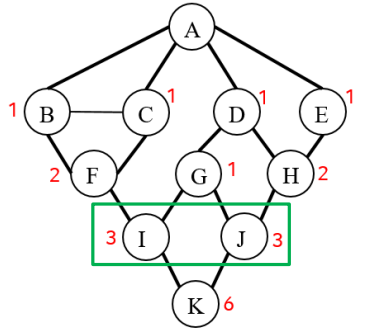
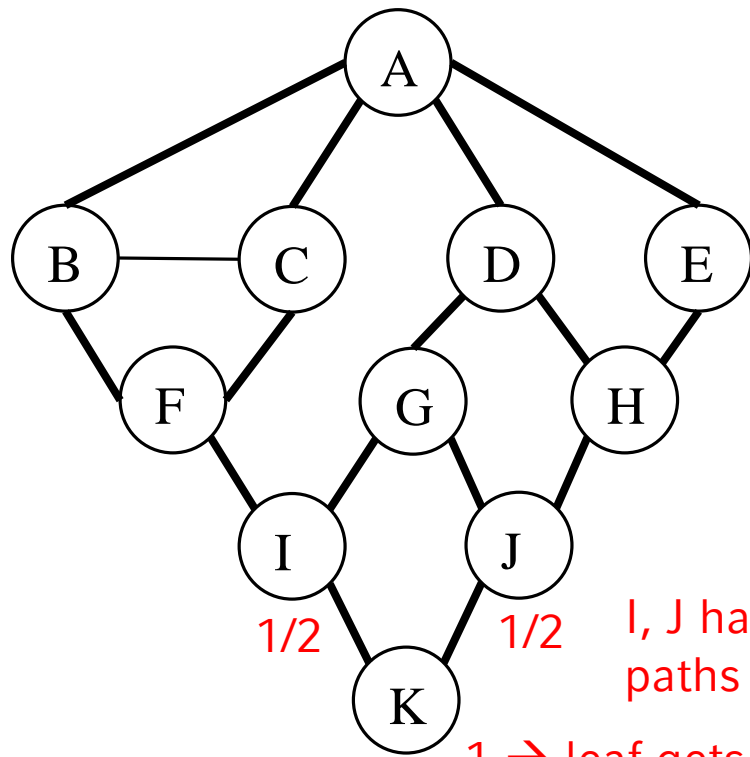
Step 3:



1 → leaf gets a credit of 1.

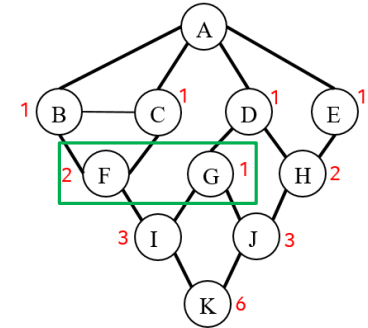
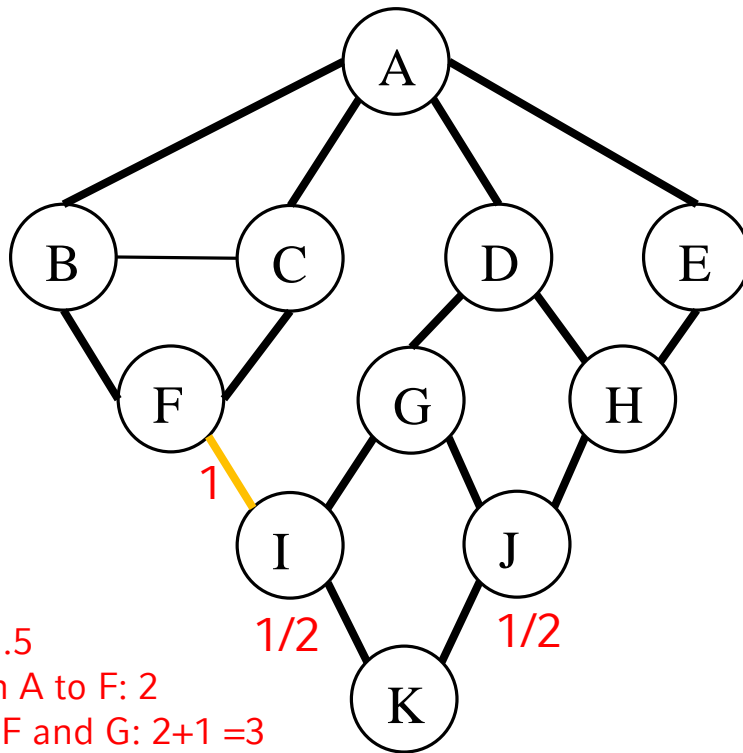


Step 3:



I, J have both equally 3 shortest paths $\rightarrow 3/(3+3) = 1/2$
 1 \rightarrow leaf gets a credit of 1.

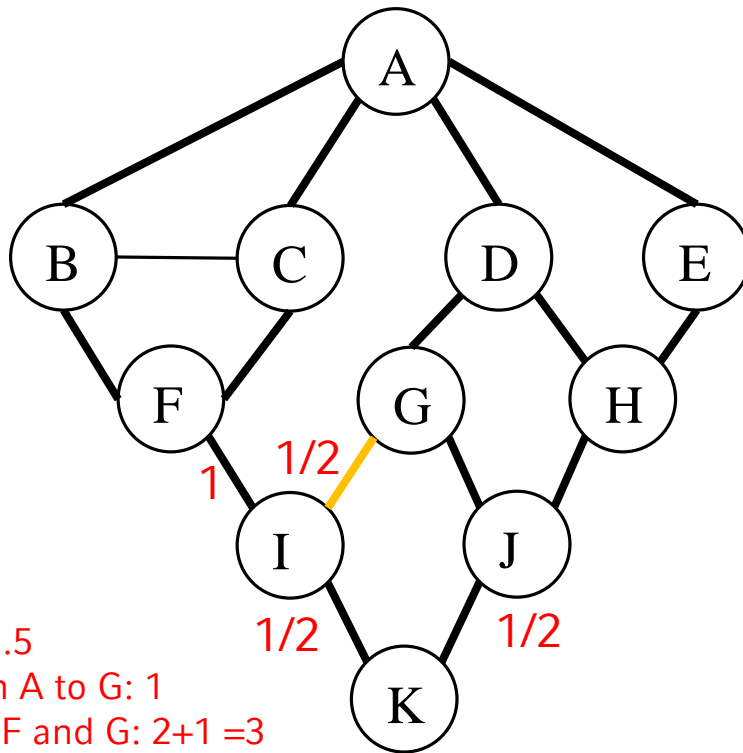
Step 3:



Credit c of I: $1+0.5=1.5$
 # shortest paths from A to F: 2
 sum #shortest paths F and G: $2+1=3$

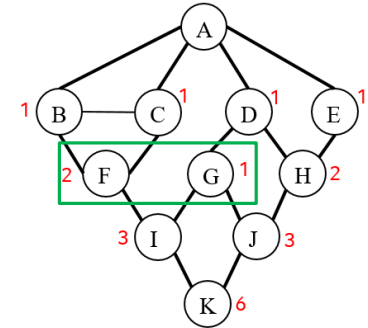
betweenness $B(I, F) = \frac{(1.5 * 2)}{3} = 1$

Step 3:



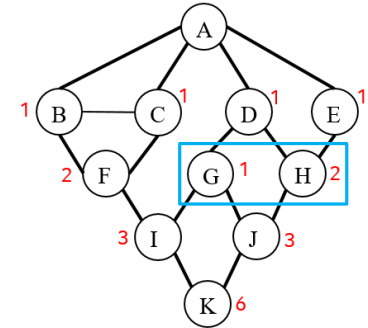
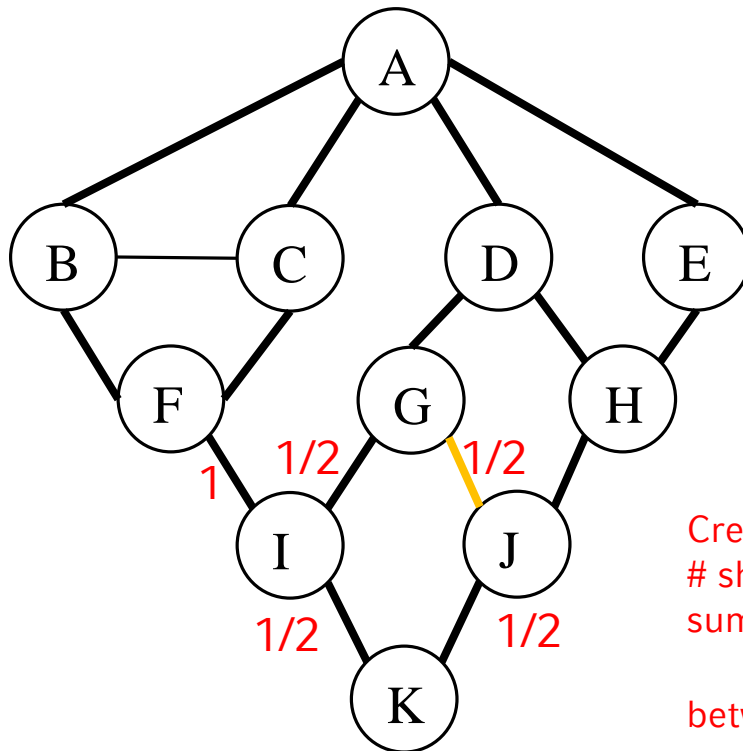
Credit c of I: $1+0.5=1.5$
 # shortest paths from A to G: 1
 sum #shortest paths F and G: $2+1 = 3$

betweenness $B(I, G) = \frac{(1.5 * 1)}{3} = 1/2$



Assignment 12-2

Step 3:

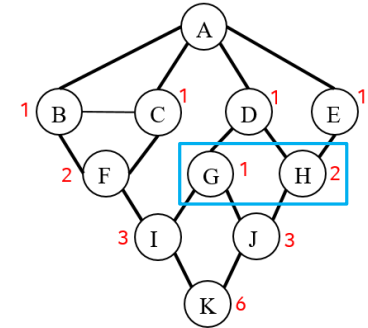
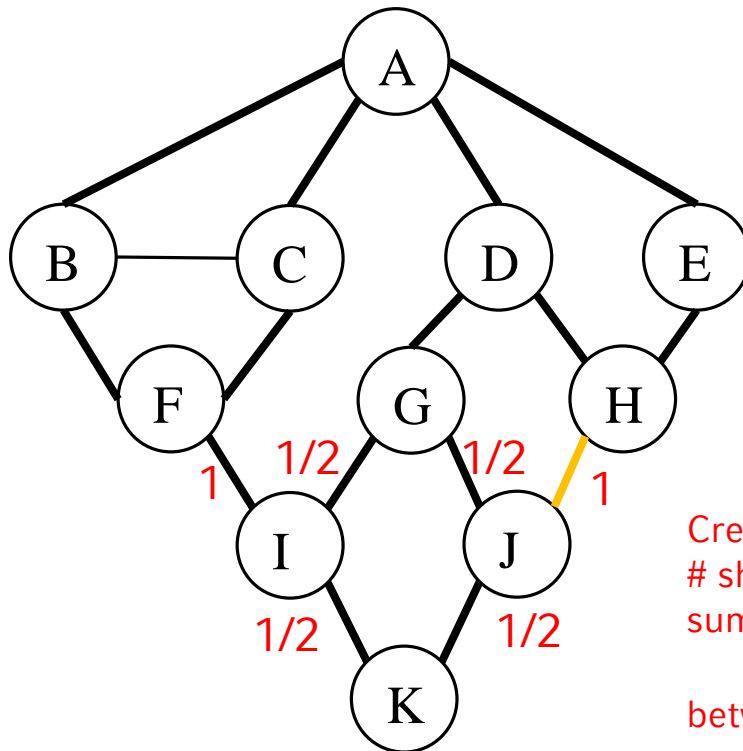


Credit c of J: $1+0.5=1.5$
 # shortest paths from A to G: 1
 sum #shortest paths G and H: $1+2=3$

betweenness $B(J, G) = \frac{(1.5 * 1)}{3} = \frac{1}{2}$

Assignment 12-2

Step 3:

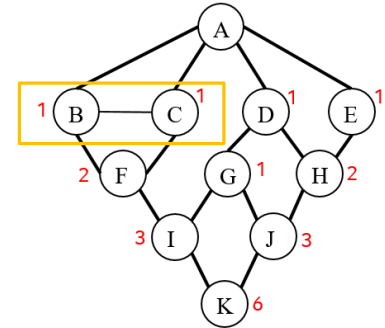
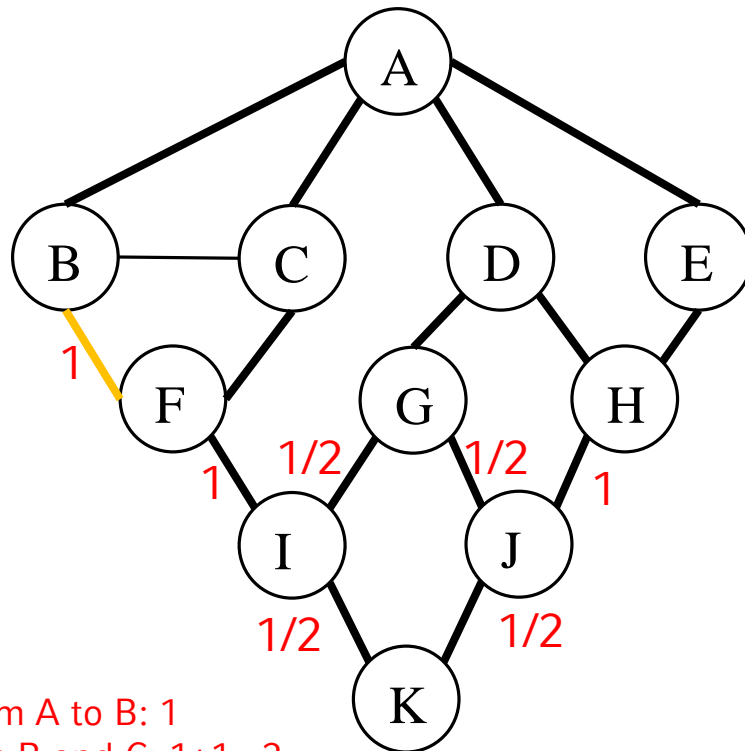


Credit c of J: $1+0.5=1.5$
 # shortest paths from A to H: 2
 sum #shortest paths G and H: $1+2=3$

betweenness $B(J, H) = \frac{(1.5 * 2)}{3} = 1$

Assignment 12-2

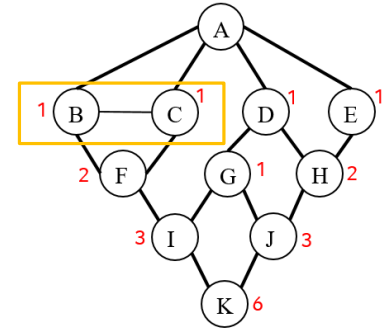
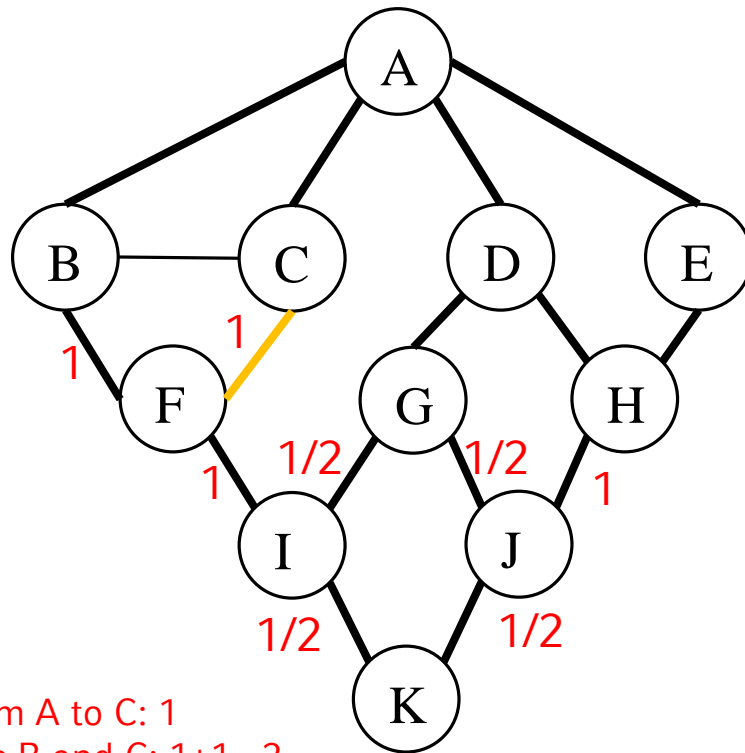
Step 3:



Credit c of F: $1+1=2$
 # shortest paths from A to B: 1
 sum #shortest paths B and C: $1+1=2$

betweenness $B(F, B) = \frac{(2*1)}{2} = 1$

Step 3:

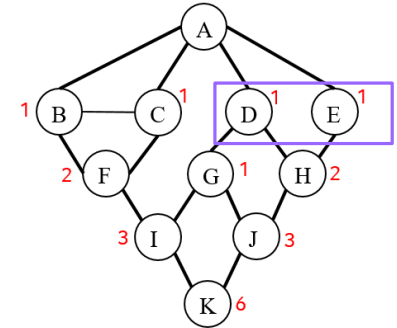
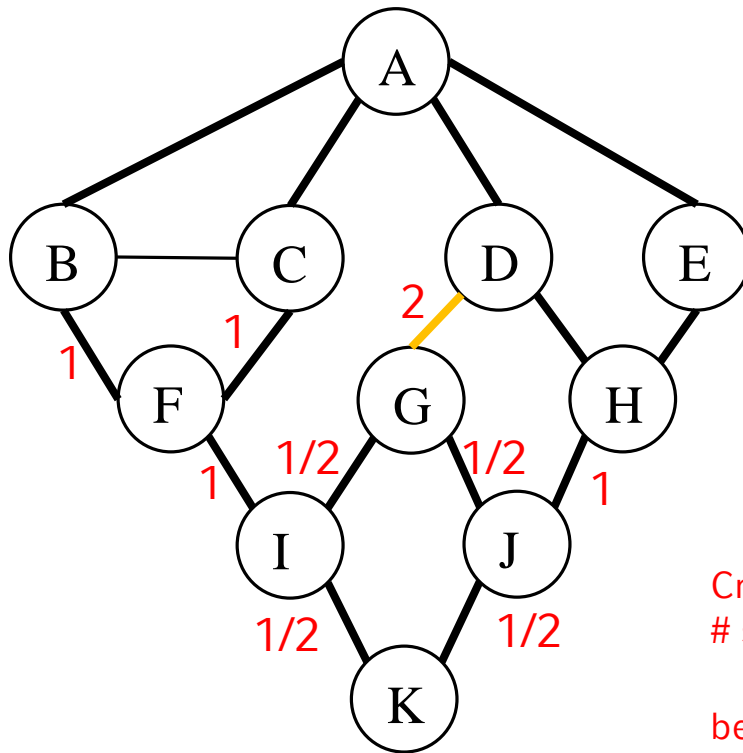


Credit c of F: $1+1=2$
 # shortest paths from A to C: 1
 sum #shortest paths B and C: $1+1=2$

betweenness $B(F, C) = \frac{(2 \cdot 1)}{2} = 1$

Assignment 12-2

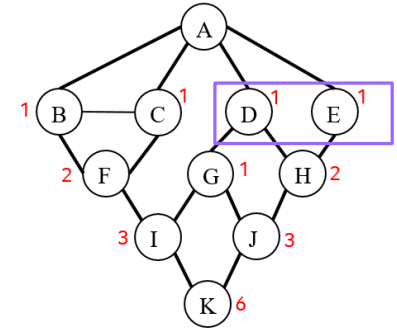
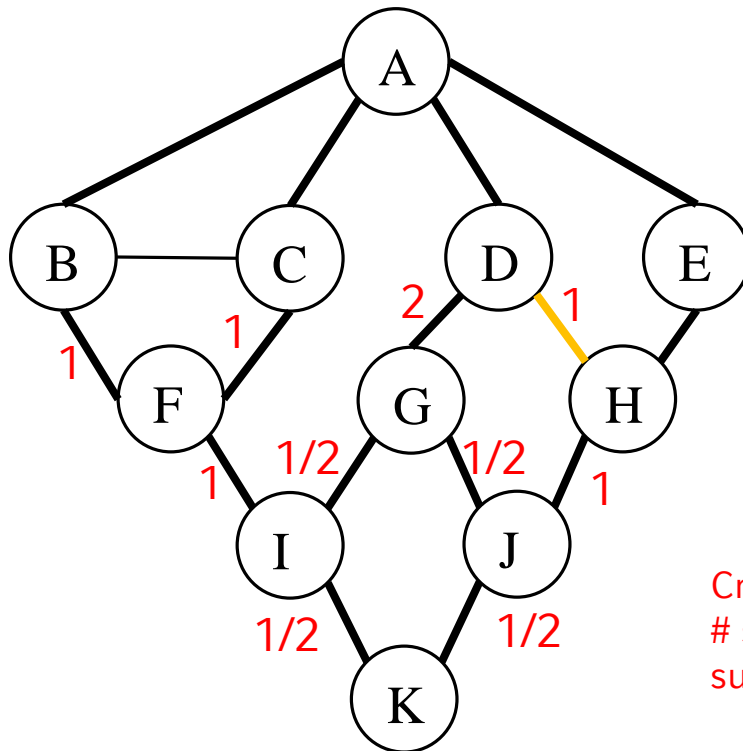
Step 3:



Credit c of G : $1+0.5+0.5=2$
 # shortest paths from A to D : 1

betweenness $B(G, D) = \frac{(2*1)}{1} = 2$

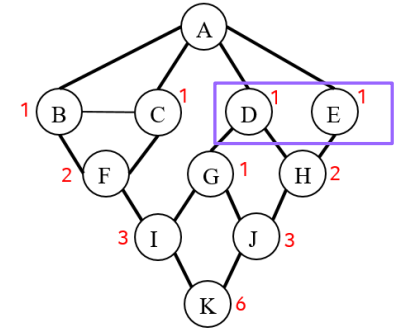
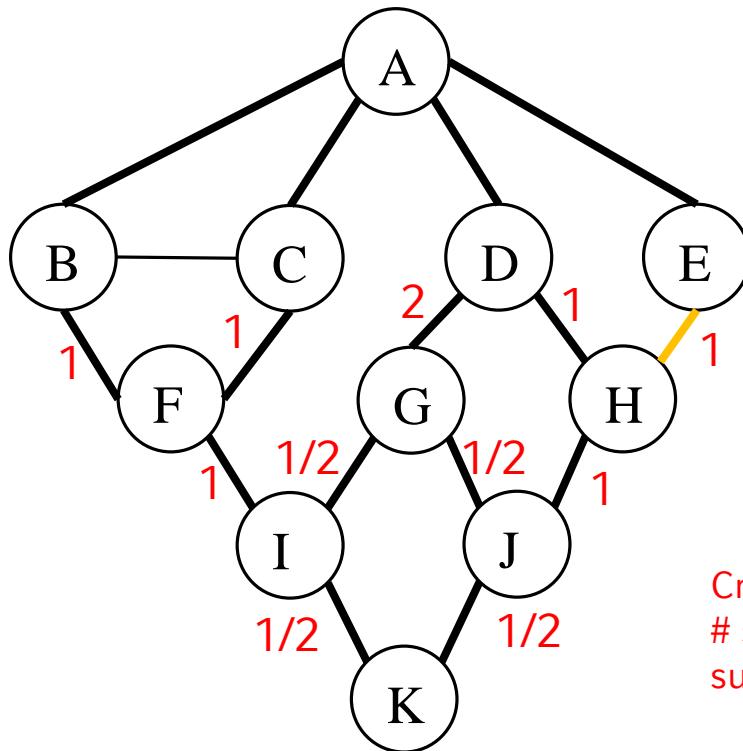
Step 3:



Credit c of H: $1+1=2$
 # shortest paths from A to D: 1
 sum #shortest paths D and E: $1+1 = 2$

betweenness $B(H, D) = \frac{(2*1)}{2} = 1$

Step 3:

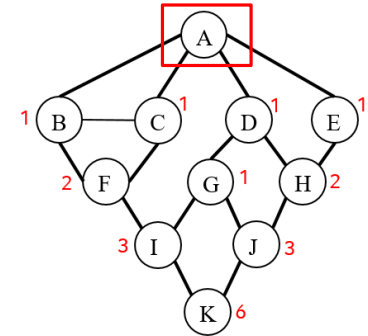
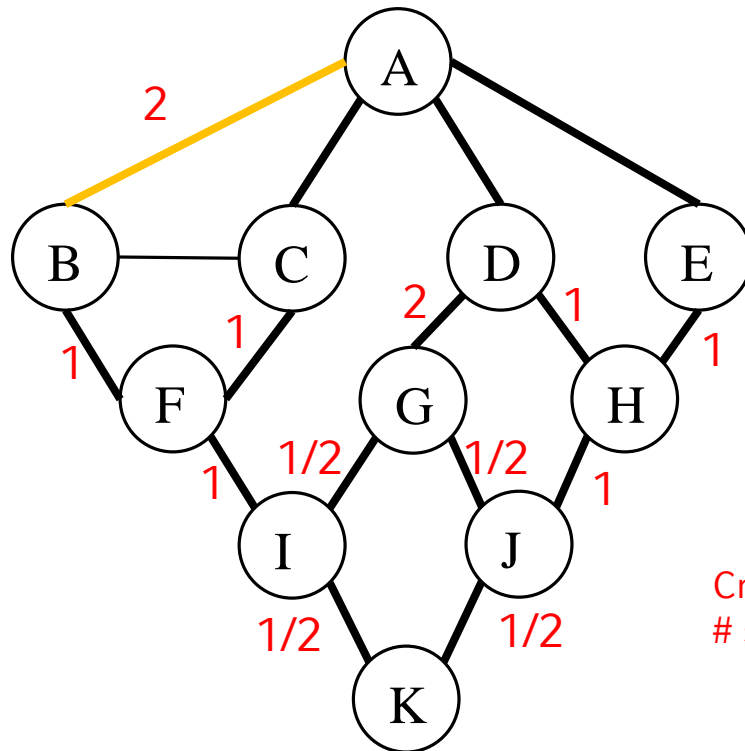


Credit c of H: $1+1=2$
 # shortest paths from A to E: 1
 sum #shortest paths D and E: $1+1=2$

betweenness $B(H, E) = \frac{(2*1)}{2} = 1$

Assignment 12-2

Step 3:

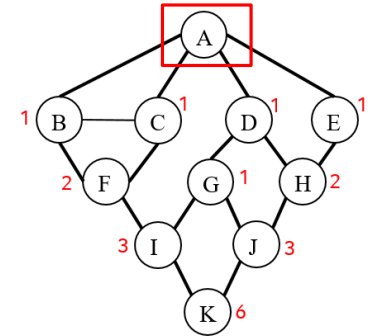
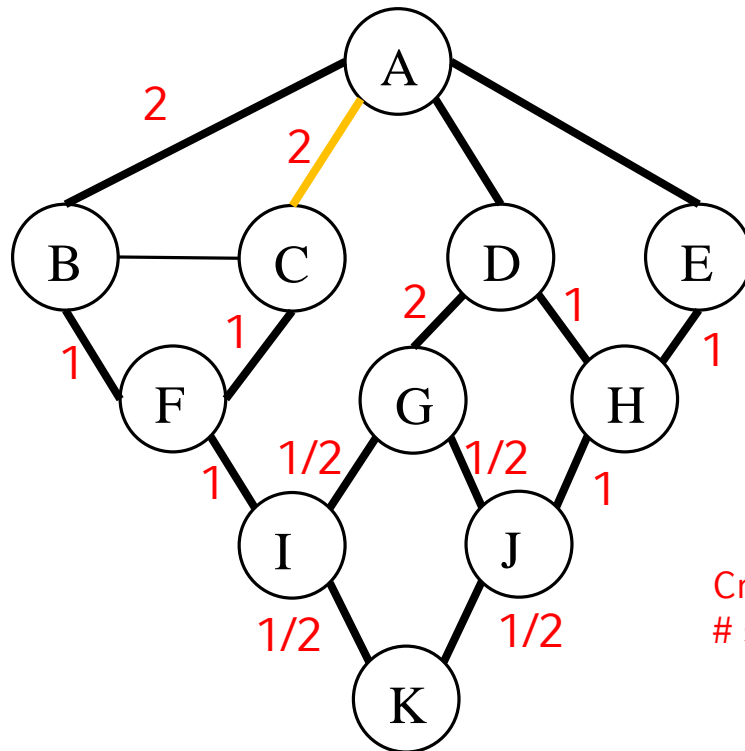


Credit c of B: $1+1=2$
 # shortest paths from A to A: 1

betweenness $B(B, A) = \frac{(2*1)}{1} = 2$

Assignment 12-2

Step 3:

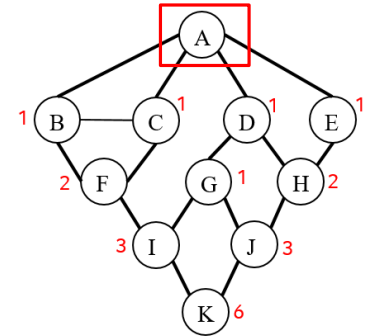
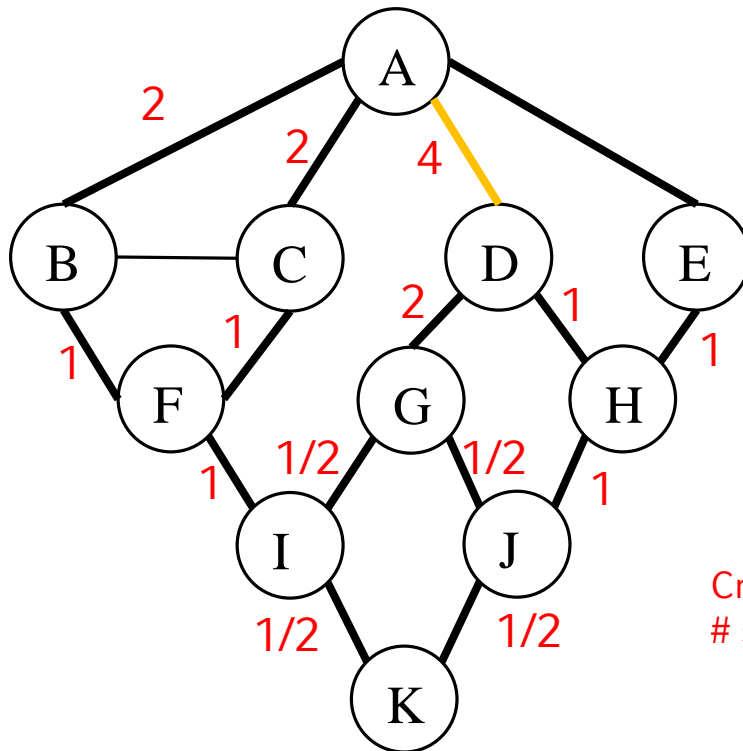


Credit c of C: $1+1=2$
 # shortest paths from A to A: 1

betweenness $B(C, A) = \frac{(2*1)}{1} = 2$

Assignment 12-2

Step 3:

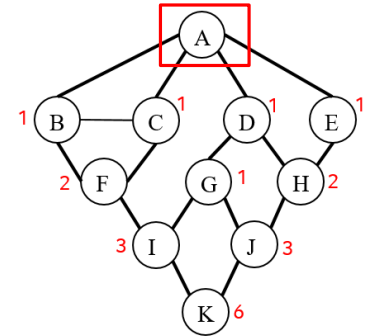
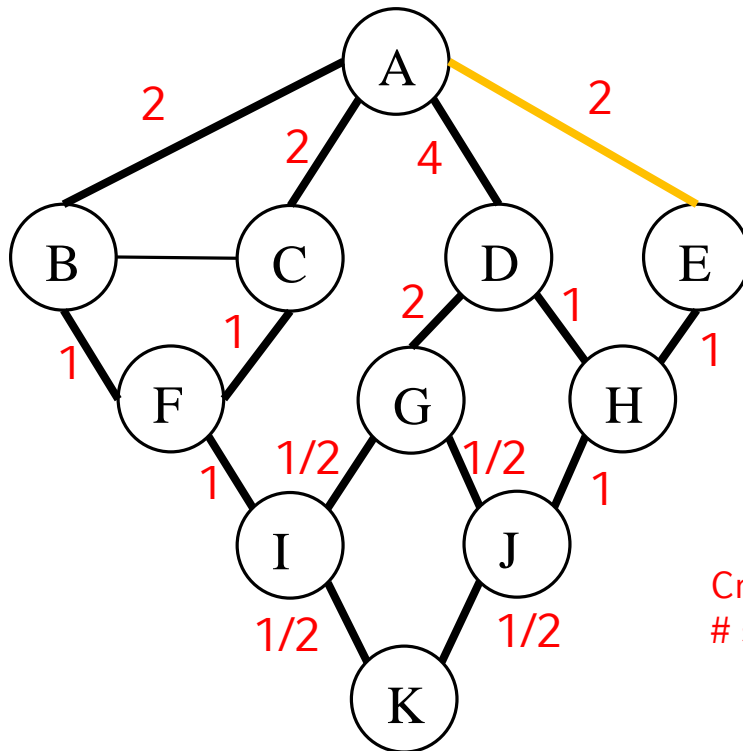


Credit c of D : $1+2+1=4$
 # shortest paths from A to A : 1

betweenness $B(D, A) = \frac{(4*1)}{1} = 4$

Assignment 12-2

Step 3:



Credit c of E : $1+1=2$
 # shortest paths from A to A : 1

betweenness $B(E, A) = \frac{(2*1)}{1} = 2$