



### **BeyOND – Unleashing BOND**

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#### 1. Background

- Motivation: k-nearest neighbor search in high-dimensional databases
- BOND revisited

#### 2. Introducing BeyOND

- Filtering objects via distance approximations
- Sub Cubes, MBRs

#### 3. Experimental Evaluation

#### 4. Conclusions

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• Similarity search in high-dimensional space is

☺ important in cases of images, e-commerce, etc.⊗ slow

- The suitability of index-based solutions depends on the data distribution
- Open question: relevant vs. irrelevant attributes
- Similarity search in subspaces:
  - Fix query attributes beforehand
  - Use multiple pivot points to derive upper and lower bounds
  - Process data vertically to reduce the high-dimensional space



## **BOND Revisited (1)**



- BOND<sup>[1]</sup>: k-nearest neighbor search on high-dimensional data
  - Resolves feature vectors (FVs) column-wise
  - Ranking of columns w.r.t. relevance
  - Pruning of columns using a branch-and-bound approach
  - Resolved part is known exactly
  - Unresolved part has to be approximated
  - Resolving stops when approximation is "good enough"
  - Support of subspace queries
  - Distance metrics:
    - Histogram intersection (uncorrelated dimensions)
    - Euclidean distance

[1] de Vries, Mamoulis, Nes, Kersten: *Efficient k-NN Search On Vertically Decomposed Data* (SIGMOD'02)





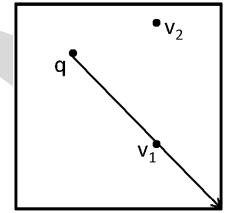
- Restrictions of BOND:
  - 1. The approach works only on Zipfian distributed data.
  - 2. The feature values are normalized to [0,1] in each dimension.
  - The proposed bounds are loose. The validity of stricter bounds (Bond advanced) depends on a certain resolve order of the columns.



- Notation:
  - query vector q, database vector v
  - Splitting of v: resolved part  $v^-$ , unresolved part  $v^+ \Rightarrow v = v^- \cup v^+$
- Approximated distance:  $S_{approx}(q,v) = S_1(q^-,v^-) + S_2(q^+,v^+)$ 
  - Resolved part:  $S_1(q^-, v^-) = \sum_{i=1}^{\infty} (q_i^- v_i^-)^2$
  - Unresolved part:  $S_2(q^+, v^+) = \sum_{i=1}^{i} \max\{q_i^+, 1 q_i^+\}^2 \ge S_1(q^+, v^+)$
- Distance bounds:

$$S_{upper}(q,v) = S_1(q^-,v^-) + S_2(q^+,v^+) \ge S_1(q,v)$$

$$S_{lower}(q, v) = S_1(q^-, v^-) + 0 \le S_1(q, v)$$







- Benefits of BeyOND:
  - 1. Independence of the data distribution.
  - 2. No restriction to a normalized data space.
  - 3. No specific resolve order of the dimensions is needed.  $\odot$

 $\Rightarrow$  Price: Distance approximations are no more suitable!  $\otimes$ 

- Solution: Combining the idea of BOND with well-known techniques:
  - VA-file (data space partitioning)
  - MBR (Minimum Bounding Rectangle) approximation (data organizing)
- ⇒ Remaining restriction: minimum/maximum values for each dimension need to be known ⊗

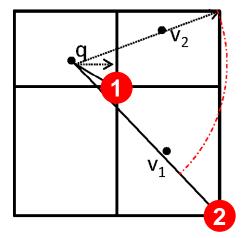


### Sub Cubes (1)



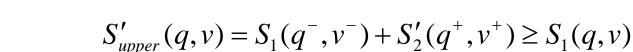
- First extension: VA-file<sup>[2]</sup> with one split
  - $\Rightarrow 2^d$  sub cubes
  - $\Rightarrow$  Addressing via Z-IDs

 $\Rightarrow$  Improved bounds based on the close / far sub cube borders  $c_{v_i}^{lower}$  1 and  $c_{v_i}^{upper}$  2



- Memory-efficient representation (8 bytes  $\rightarrow$  1 bit)
  - Sub cube need not be kept in main memory
- Split positions stored in one separate array per dimension
- Dependence on split level:
  - FV: 8 bytes per dimension
  - s splits: s / 8 bytes (s bits) per dimension

[2] Weber, Schek, Blott. A Quantitative Analysis and Performance Study for Similarity Search Methods in High-Dimensional Spaces (VLDB'98)



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New distance bounds:



• Old distance bounds:

$$S_{upper}(q, v) = S_1(q^-, v^-) + \sum_i \max\{q_i^+, 1 - q_i^+\}^2 - S_{lower}(q, v) = S_1(q^-, v^-) + 0$$

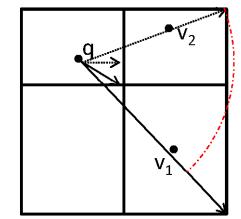
• Approximations of unresolved dimensions:

$$S_{2}'(q^{+},v^{+}) = \sum_{i} \max \left\{ q_{i}^{+} - c_{v_{i}^{+}}^{lower} \Big|, \Big| q_{i}^{+} - c_{v_{i}^{+}}^{upper} \Big| \right\}^{2}$$

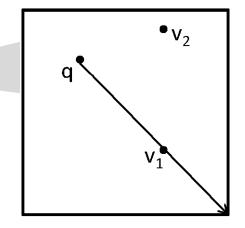
$$S_{2}''(q^{+},v^{+}) = \sum_{i} \left\{ \begin{array}{c} 0 & \text{if } q_{i}^{+} \\ \min \left\{ q_{i}^{+} - c_{v_{i}^{+}}^{lower} \Big|, \Big| q_{i}^{+} - c_{v_{i}^{+}}^{upper} \Big| \right\}^{2} & else \end{array}$$

 $S'_{lower}(q,v) = S_1(q^-,v^-) + S''_2(q^+,v^+) \le S_1(q,v)$ 

if  $q_i^+ \in \left[c_{v_i^+}^{lower}, c_{v_i^+}^{upper}\right]$ 







- 8 byte coordinates: memory increase is limited to  $\frac{d \cdot 16}{card(MBR)}$ per feature vector (+ pointer to Z-ID)

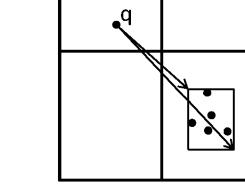
 $f(MBR) = \frac{V_{subcube}}{V_{MBR}} \cdot card(MBR)$ 

bytes

Dense sub cubes allow a tighter approximation via MBRs

**MBR Caching (1)** 

- Restrict the number of MBRs in order to avoid a memory overhead
- Ranking function for MBRs:





Most sub cubes are (very) sparse, i.e. occupied by at most





•

one FV



### **MBR Caching (2)**



- Limit the number of MBRs to 1% of the database size
- Threshold as a trade-off between pruning power and additional memory consumption
- Requirements:
  - Either all MBRs can be kept in memory,
  - or the time for loading the MBRs is less than the time for resolving the respective FVs.
- Adaption of the equations for lower and upper bounds





- Evaluated approaches:
  - 1. BondAdvanced (stricter bounds, but resolve order dependent)
  - 2. Bond (original bounds)\*
  - 3. Sequential\*
  - 4. Beyond-1 (1 split)
  - 5. BeyondMBR-1 (1 split + MBRs)
  - 6. Beyond-2
  - 7. BeyondMBR-2
  - 8. Beyond-3\*
  - 9. BeyondMBR-3\*





#### • Data set descriptions:

Data Set	Dims	Size	Туре
ALOI	27	110,250	Color Histograms, Zipfian
CLUSTERED	20	500,000	Synthetic, 50 Clusters, Gaussian
PHOG <sup>[3]</sup>	110	10,715	CT Histograms, PCA'ed
SIFT <sup>[4]</sup>	133	335,583	Image Features

[3] Graf, Kriegel, Schubert, Poelsterl, Cavallaro. 2D Image Registration in CT Images Using Radial Image Descriptors (MICCAI'11)

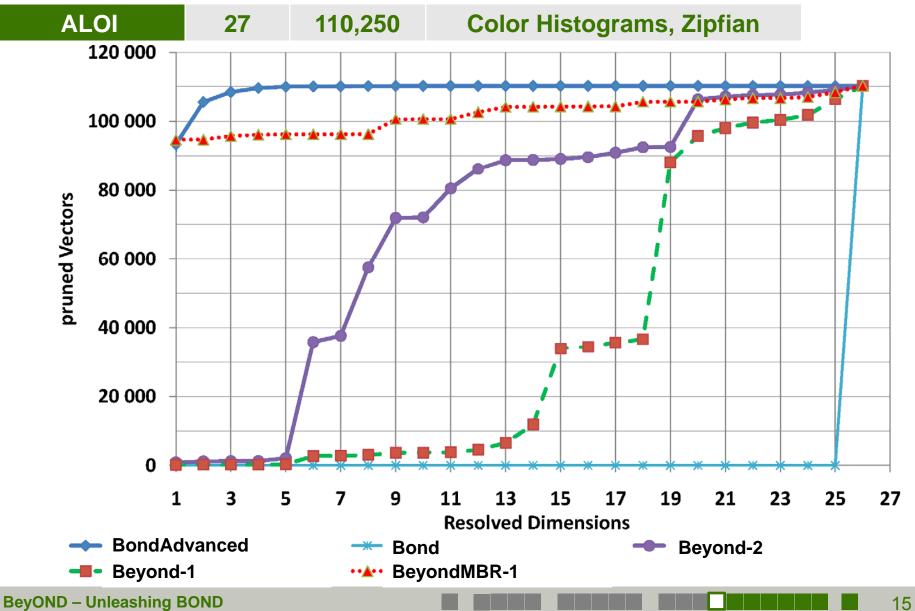
[4] Lowe. Distinctive Image Features from Scale-Invariant Keypoints (Int. Journal of Computer Vision, 2004)

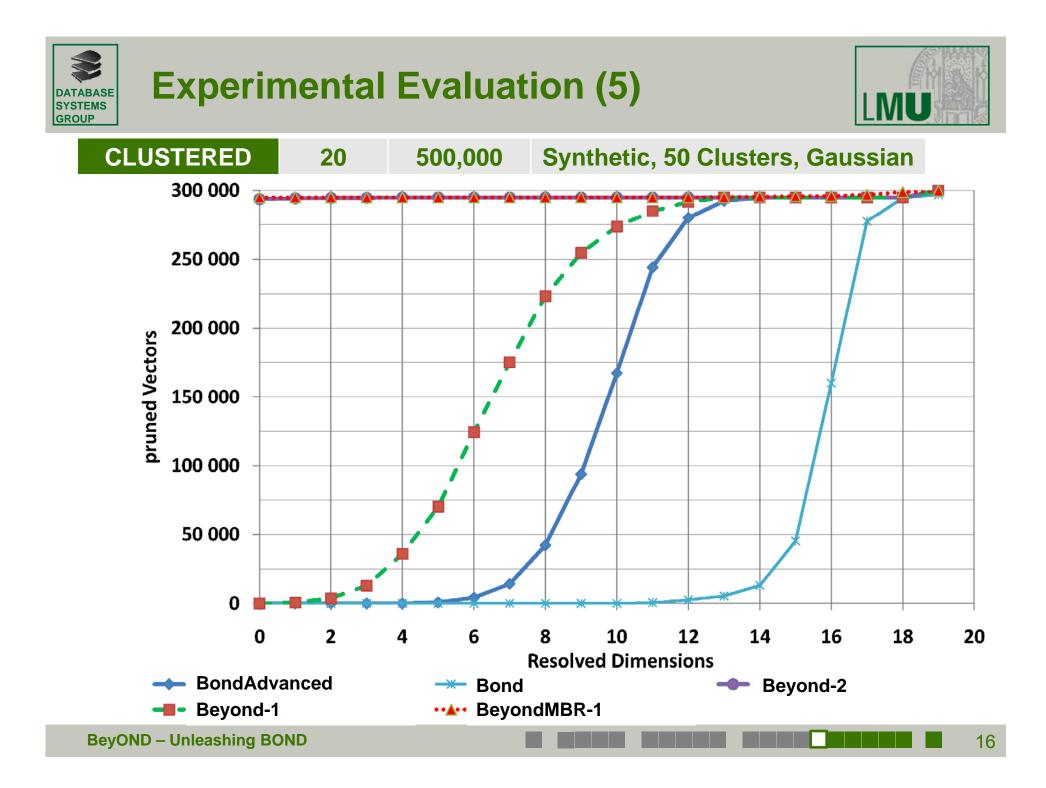




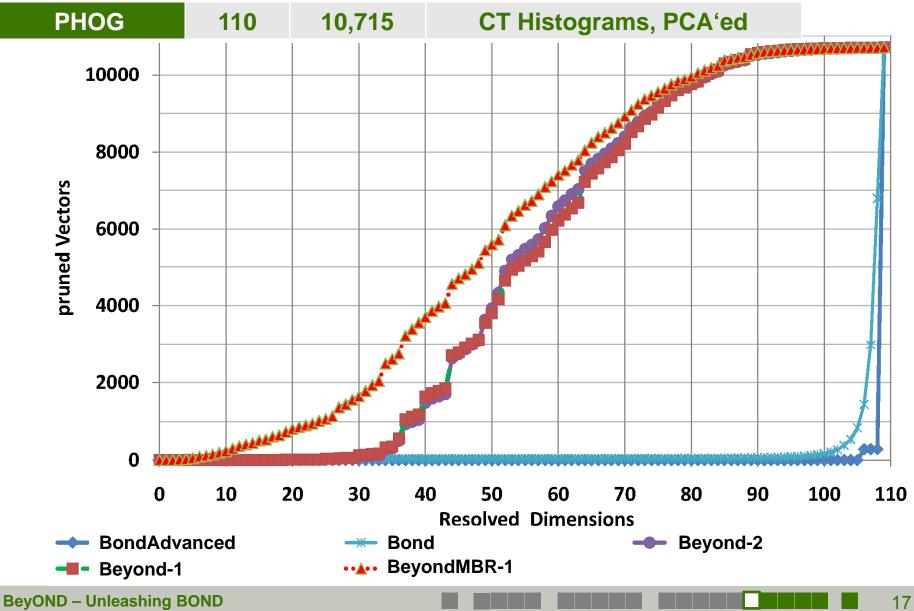
- Experimental settings:
  - 50 k-nearest neighbor queries
  - k = 10
  - Averaged cumulative number of pruned FVs after resolving a column
  - AUC: data not resolved
  - AOC: data resolved for refinement













### **Experimental Evaluation (7)**



#### Pruning power (Sub cubes)

Data Set	Splits	25% pruned	50% pruned	90% pruned
ALOI	1	16 (59%)	19 (70%)	23 (85%)
CLUSTERED	1	7 (35%)	8 (40%)	10 (50%)
PHOG	1	45 (41%)	58 (53%)	80 (73%)
ALOI	2	7 (26%)	9 (33%)	21 (75%)
CLUSTERED	2	1 (5%)	1 (5%)	1 (5%)
PHOG	2	45 (41%)	55 (50%)	79 (72%)

Pruning power	Data Set	Splits	25% pruned	50% pruned	90% pruned
(Sub cubes +	ALOI	1	1 (4%)	1 (4%)	10 (37%)
MBRs)	CLUSTERED	1	1 (5%)	1 (5%)	1 (5%)
,	PHOG	1	37 (34%)	50 (45%)	77 (70%)

#	# Accessed				
	columns				

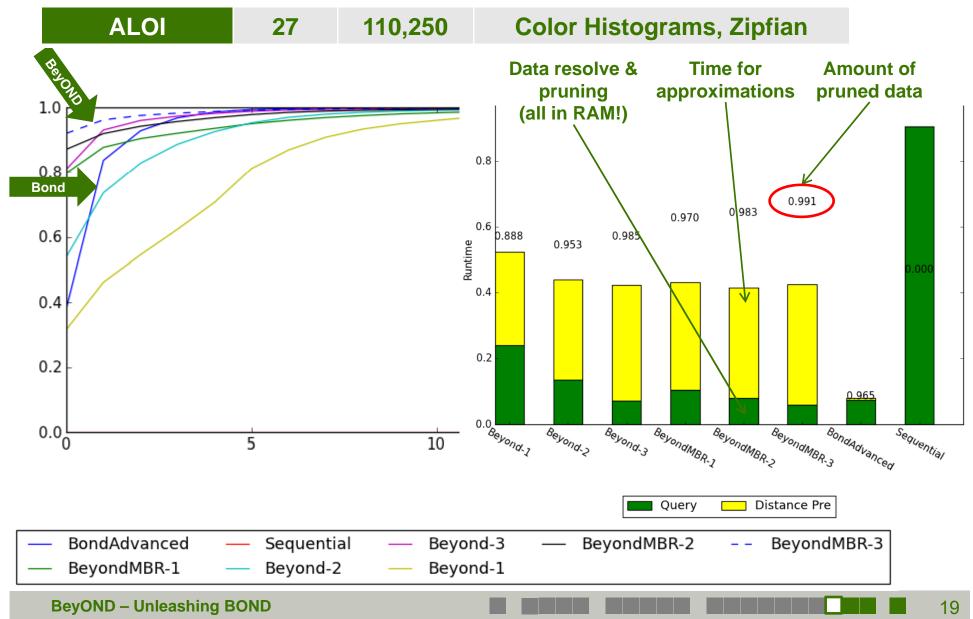
essed	Data Set	1 split	2 splits	1 split + MBR
umns	ALOI	66.9%	38.3%	7.7%
	CLUSTERED	34.1%	1.6%	1.4%
	PHOG	52.6%	52.3%	45.4%

## **Experimental Evaluation (8)**

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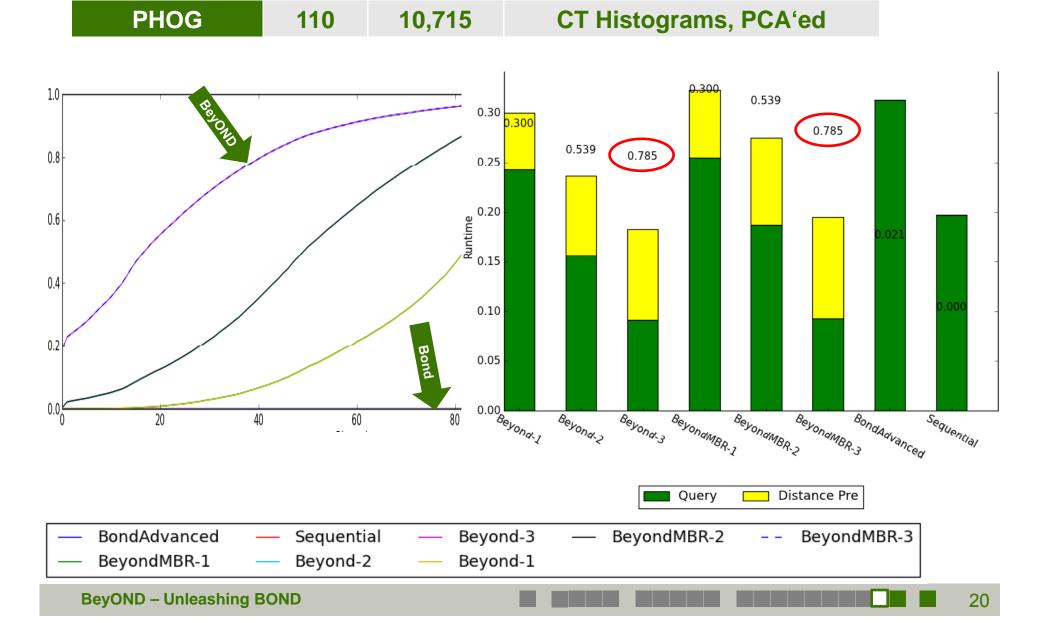


#### **Experimental Evaluation (9)** DATABASE

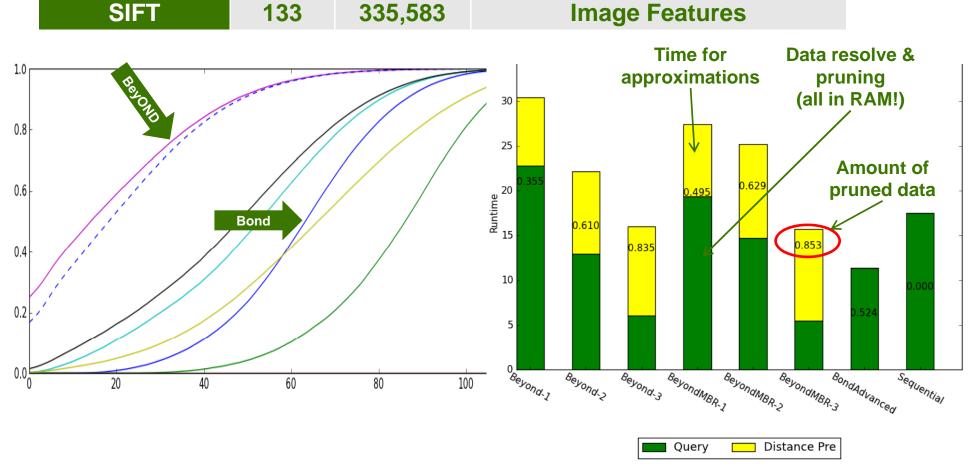
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BondAdvanced
 BeyondAdvanced
 Beyond-2
 Beyond-3
 BeyondMBR-2
 BeyondMBR-2
 BeyondMBR-2
 BeyondMBR-3





- Removed restrictions...
  - 1. Independence of the data distribution.
  - 2. No restriction to a normalized data space.
  - 3. No specific resolve order of the dimensions is needed.
- Combination of relevant techniques...
  - VA-file-based partitioning of the data space
  - MBR caching
- Still open issues...
  - Trade-off: split level vs. pruning power
  - Trade-off: MBR memory consumption vs. pruning power
  - Sophisticated techniques for the creation of the MBRs
  - Overcome the restriction that the vector lengths have to be known





# Thank you for listening!

# Any questions?

http://www.dbs.ifi.lmu.de/cms/Publications/BeyOND\_-\_Unleashing\_BOND