

## SCAN: A Structural Clustering Algorithm for Networks

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#### scaling: #edges connected to respective vertex

## **Network basics**

- Network graph  $G = \{V, E\}$
- *V* is set of vertices, *E* is set of edges connecting the vertices

- Hubs: vertices that bridge (many) clusters
- **Outliers**: vertices marginally connected to clusters featuring a weak association
- network clustering/graph partitioning: division of graph into set of disjoint sub-graphs in order to find hidden structures

Hub

Outlier

• existing approaches aim at finding clusters based on the number of edges between vertices or clusters

## SCAN – A Structural Clustering Algorithm for Networks



- based on definitions of DBSCAN
- detects clusters, hubs and outliers using structure and connectivity of the vertices as a clustering criterion
- vertices sharing a certain quantity of neighbors should be grouped into one cluster, hubs and outliers should be isolated
- runtime complexity wrt. *n* vertices and *m* edges in a network: *O*(*m*)
- focusses on simple, undirected and unweighted graphs

## **Definitions for structure-connected clusters (I)**



VERTEX STRUCTURE  $v \in V$ , the structure of v is defined by its neighborhood  $\Gamma(v)$  $\Gamma(v) = \{ w \in V | (v, w) \in E \} \cup \{ v \}$ 8 12 **STRUCTURAL SIMILARITY** 13  $\sigma(4,5) = \frac{4}{\sqrt{5 \cdot 6}} = 0,73$ • two objects sharing a similar structure in terms of their neighborbood, will have a large structural similarity  $\sigma(v,w) = \left\{ \frac{|\Gamma(v) \cap \Gamma(w)|}{\sqrt{|\Gamma(v)||\Gamma(w)|}} \right\}$ 

## Definitions for structure-connected clusters (II)



#### $\varepsilon$ – NEIGHBORHOOD

apply threshold to structural similarity in order to assign cluster memberships

$$N_{\varepsilon}(v) = \{ w \in \Gamma(v) \mid \sigma(v, w) \ge \varepsilon \}$$



• CORE VERTEX

core vertex shares *structural similiarity* of at least  $\epsilon$  with at least  $\mu$  neighbors

 $Core_{\varepsilon,\mu}(v) \Leftrightarrow |N_{\varepsilon}(v)| \geq \mu$ 

## Definitions for structure-connected clusters (III)



• **DIRECT STRUCTURE REACHABILITY** if vertex is in  $\varepsilon$  – Neighborhood of a *core vertex*, they should be in the same cluster  $DirReach_{\varepsilon,\mu}(v,w) \Leftrightarrow Core_{\varepsilon,\mu}(v) \land w \in N_{\varepsilon}(v)$ 



#### • STRUCTURE REACHABILITY

a vertex  $w \in V$  is *structure-reachable* from  $v \in V$ , if there is a chain of vertices  $v_1, ..., v_n \in V, v_1 = v, v_n = w$  such that  $v_{i+1}$  is *directly structure-reachable* from  $v_i$   $Reach_{\varepsilon,\mu}(v,w) \Leftrightarrow \exists v_1, ..., v_n \in V: v_1 = v \land v_n = w \land \forall i$   $\in \{1, ..., n\}: DirReach_{\varepsilon,\mu}(v_i, v_{i+1})$  $\varepsilon = 0.65, \mu = 2$ 

## Definitions for structure-connected clusters (IV)



#### • STRUCTURE CONNECTIVITY

two non-core vertices  $v, w \in V$  belonging to the same cluster may not be *structure-reachable*, as the *core-condition* does not hold for them.

However, they still belong to the same cluster, if they are *structure-reachable* from the same *core vertex*  $u \in V$ . This is called *structure-connectivity*.

 $Connect_{\varepsilon,\mu}(v,w) \Leftrightarrow \exists u \in V: Reach_{\varepsilon,\mu}(u,v) \land Reach_{\varepsilon,\mu}(u,w)$ 



## Definitions for structure-connected clusters (V)



#### • STRUCTURE-CONNECTED CLUSTER

a non-empty subset  $C \subseteq V$  is called a *structure-connected cluster*, if all vertices in C are *structure-connected* and C is maximal wrt. *structure reachability* 

 $\begin{array}{l} Cluster_{\varepsilon,\mu}(\mathcal{C}) \Leftrightarrow \\ 1. \ \text{Connectivity:} \ \forall v, w \in \mathcal{C}: \ Connect_{\varepsilon,\mu}(v,w) \\ 2. \ \text{Maximality:} \ \forall v, w \in V: v \in \mathcal{C} \land Reach_{\varepsilon,\mu}(v,w) \Rightarrow w \in \mathcal{C} \end{array}$ 

#### • CLUSTERING

a clustering P consists of all structure-connected clusters

 $Clustering_{\varepsilon,\mu}(P) \Leftrightarrow P = \{C \subseteq V \mid Cluster_{\varepsilon,\mu}(C)\}$ 

## **Definitions for structure-connected clusters (VI)**



If a vertex is not a member of any cluster, it is either a hub or an outlier, depending on its neighborhood.

• HUB

if an isolated vertex  $v \in V$  has neighbors belonging to two or more different clusters, it is a *hub* 

 $Hub_{\varepsilon,\mu}(v) \Leftrightarrow$ 

1. *v* is not a member of any cluster:  $\forall C \in P : v \notin C$ 

2. *v* bridges different clusters:

 $\exists p, q \in \Gamma(v) : \exists X, Y \in P : X \neq Y \land p \in X \land q \in Y$ 

#### • OUTLIER

an isolated vertex  $v \in V$  is an *outlier*, if and only if all its neighbors either belong to only one cluster or do not belong to any cluster

 $Outlier_{\varepsilon,\mu}(v) \Leftrightarrow$ 

1. *v* is not a member of any cluster:  $\forall C \in P : v \notin C$ 

2. *v* does not bridge different clusters:  $\neg \exists p, q \in \Gamma(v)$ :  $\exists X, Y \in P$ :  $X \neq Y \land p \in X \land q \in Y$ 



## **Algorithm SCAN**



- for each vertex **v** not yet classified, it is checked whether vertex is a **core** 
  - **if so**, a new cluster is expanded from this vertex:
  - 1. the algorithm generates a **new cluster id**, looks for **all unclassified vertices** in the **neighborhood** of the core and inserts them into a **queue**

vertices which are directly structure-reachable from core vertex v

2. it traverses the queue until it is empty and establishes all vertices which are **directly structure-reachable** from the respective vertex **w** 

vertices which are structure-reachable from core vertex v (via vertex w)

- 3. the same cluster id is assigned to all those vertices and if they are not labeled as a non-member yet, they are also inserted into the queue
- if not, it is labeled as a non-member
- non-member vertices which have edges to two or more clusters, are classified as hubs. Otherwise, they are classified as outliers.
- $\epsilon$ -value between 0.5 and 0.8,  $\mu$  = 2

## **Experimental Evaluation (I)**



Comparison to **FastModularity** algorithm:

- hierarchical network clustering algorithm optimizing the modularity
- *modularity* measures whether division of a network into communities is a good one in terms of many edges within communities and preferably little edges between communities
- 1. initally, each vertex is the only member of a community
- 2. iteratively, the algorithm greedily merges the two communities causing the largest increase of modularity, until all vertices are members of the same community

### **Experimental Evaluation (II)** Customer Segmentation dataset





## **Experimental Evaluation (III)**





SCAN clustering result



### **Experimental Evaluation (IV)** Books about US politics dataset







### Discussion



- + identifies not only clusters, but also outliers and hubs
- + linear run-time complexity wrt. # edges, each vertex is visited only once
- performance highly depends on sensitive input parameters
- ignores domain knowledge in clustering attributes
- assumes that network is homogeneous and adjacency matrix is already defined



# Q&A V