FAKULTÄT FÜR MATHEMATIK, INFORMATIK UND STATISTIK INSTITUT FÜR INFORMATIK

LEHRSTUHL FÜR DATENBANKSYSTEME UND DATA MINING

Lecture Notes to Big Data Management and Analytics Winter Term 2018/2019 NoSQL Databases

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NoSQL Database Systems

Outline

- History
- Concepts
 - ACID
 - BASE
 - CAP
- Data Models
 - Key-Value Stores
 - Document Databases
 - Wide Column Stores
 - Graph Databases

60s: IBM developed the Hierarchical Database Model

- Tree-like structure
- Data stored as *records* connected by *links*
- Support only one-to-one and one-to-many relationships

Mid 80's: Rise of Relational Database Model

- Data stored in a collection of tables (rows and columns)
 → Strict relational scheme
- SQL became standard language (based on relational algebra)
- → Impedance Mismatch!

History – Impedance Mismatch

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SNR: L1									
Sname: Meier status: 20									
location: Wetter									
Project: PNR· P2									
Pname: Pleite									
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ONR: T6					- 7				
Uname: screw									
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Given the following database scheme and an object of type Supply:

How to incorporate the data bundled in the object Supply into the relational DB?

History – Impedance Mismatch



INSERT INTO L VALUES (Supply.getSupplier().getLNR(), ...);

INSERT INTO P VALUES (Supply.getProject().getPNR(), ...);

. . .

History – Impedance Mismatch



INSERT INTO LTP VALUES (...);

- Object-oriented encapsulation vs. storing data distributed among several tables
 - \rightarrow Lots of data type maintenance by the programmer

Mid 90's: Trend of the Object-Relational Database Model

- data stored as objects (including data and methods)
- avoids of object-relational mapping
 → Programmer-friendly
- but still Relational Databases prevailed in the 90's

Mid 2000's: Rise of Web 2.0

- Lots of user generated data through web applications
 - \rightarrow storage systems had to become scaled up

Approaches to scale up storage systems

- Two opportunities to solve the rising storage system:
 - Vertical scaling Enlarge a single machine
 - Limited in space
 - Expensive
 - Horizontal scaling
 Use many commodity machines and form computer clusters or grids
 - Cluster maintenance

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Expensive

- Horizontal scaling
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Mid 2000's: Birth of the NoSQL Movement

- Problem of computer clusters: Relational databases do not scale well horizontally
- → Big Players like Google or Amazon developed their own storage systems: NoSQL ("Not-Only SQL") databases were born

Today: Age of NoSQL

Several different NoSQL systems available (>225)



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Characterstics of NoSQL Databases

There is no unique definition but some characteristics for NoSQL Databases:

- Horizontal scalability (cluster-friendliness)
- Non-relational
- Distributed
- Schema-less
- Open-source (at least most of the systems)

ACID – The holy grail of RDBMSs:

- <u>Atomicity:</u> Transactions happen entirely or not at all. If a transaction fails (partly), the state of the database is unchanged.
- <u>Consistency</u>: Any transaction brings the database from one valid state to another and does not break one of the predefined rules (like constraints).
- <u>Isolation</u>: Concurrent execution of transactions results in a system state that would be obtained if transactions were executed serially.
- <u>Durability</u>: Once a transaction has been commited, it will remain so.

BASE – An artificial concept for NoSQL databases:

- <u>Basically Available:</u> The system is generally available, but some data might not at any time (e.g. due to node failures)
- <u>Soft State:</u> The system's state changes over time. Stale data may expire if not refreshed.
- <u>Eventual consistency</u>: The system is consistent from time to time, but not always. Updates are propagated through the system if there is enough time.
- → BASE is settled on the opposite site to ACID when considering a "consistency-availability spectrum"



Data replication



The two types of consistency:

- Logical consistency: Data is consistent within itself (Data Integrity)
- Replication consistency: Data is consistent across multiple replicas (on multiple machines)

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Levels of Consistency:

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	Monot	tonic Read Consistency										
		M.R.(
		Immediate Consistency										
					Transactions							
				_		Road-Vour-C	Writes					
						Reau-Tour-C						

Levels of Consistency:

- <u>Eventual Consistency:</u> Write operations are not spread across all servers/partitions immediately
- <u>Monotononic Read Consistency</u>: A client who read an object once will never read an older version of this object
- <u>Read Your Own Writes</u>: A client who wrote an object will never read an older version of this object
- <u>Immediate Consistency</u>: Updates are propagated immediately, but not atomic

Levels of Consistency:

- <u>Strong consistency</u>: Updates are propagated immediately + support of atomic operations on single data entities (usually on master nodes)
- <u>Transactions</u>: Full support of ACID transaction model

Brewer's CAP Theorem:



Any networked shared-data system can have at most two of the three desired properties!

Possible DB-Systems by CAP Theorem:

- <u>CP-Systems:</u> Fully consistent and partitioned systems renounce availability. Only consistent nodes are available.
- <u>AP-Systems:</u> Fully available and partitioned systems renounce consistency. All nodes answer to queries all the time, even if answers are inconsistent.
- <u>AC-Systems</u>: Fully available and consistent systems renounce partitioning. Only possible if the system is not distributed.

Big Picture

CAP Theorem:

All clients always have the same view of the data

Each client can always read and write

The system works well despite physical network partitions

Ρ



The 4 Main NoSQL Data Models:

- Key/Value Stores
- Document Stores
- Wide Column Stores
- Graph Databases

Key/Value Stores:

- Most simple form of database systems
- Store key/value pairs and retrieve values by keys
- Values can be of arbitrary format



Key/Value Stores:

- consistency models range from eventual consistent to serializable
- some systems support ordering keys which enables efficient query processing, e.g., for range queries
- some systems support in-memory data maintenance and others manage data based on hard drives and SSDs

→ Available Systems are very heterogeneous

Key/Value Stores - Redis:



- in-memory data structure store with built-in replication, transactions and different levels of on-disk persistence
- supports complex types like lists, sets, hashes, ...
- supports various *atomic* operations

```
>> SET val 1
>> GET val => 1
>> INCR val => 2
>> LPUSH my_list a (=> 'a')
>> LPUSH my_list b (=> 'b', 'a')
>> RPUSH my_list c (=> 'b', 'a', 'c')
>> LRANGE my_list 0 1 => b,a
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```

Key/Value Stores – The Redis cluster model:



- data is automatically sharded across nodes
- some degree of availability, achieved by master-slave architecture (but cluster stops in the event of larger failures)
- easily extendable

Key/Value Stores – the redis cluster model:



Key/Value Stores – The Redis cluster model:





Hash slots 5001 – 10000 cannot be used anymore Master Slave Replicated Nodes Hash slots Nodes Hash slots 5000 5001 5001 B' ł 10000 10000 10001 10001 14522 14522

slave node B' is promoted as the new master and hash slots 5001 – 10000 are still available

Big Picture



Document Stores:

- store documents in form of XML or JSON
- semi-structured data records that do not have a homogeneous structure
- columns can have more than one value (arrays)
- documents include internal structure, or metadata
- data structure enables efficient use of indexes

Document Stores:

given following text:

Max Mustermann Musterstraße 12 D-12345 Musterstadt

<contact>

<first_name>Max</first_name> <last_name>Mustermann</last_name> <street>Musterstraße 12</street> <city>Musterstadt</city> <zip>12345</zip> <country>D</country> </contact>

\rightarrow find all <contact>s where <zip> is "12345"

Document Stores:



- data stored as documents in binary representation (BSON)
- similarly structured documents are bundled in collections
- provides own ad-hoc query language
- supports ACID transactions on document level

Document Stores:



MongoDB Data Management:

- automatic data sharding
- automatic re-balancing
- multiple sharding policies:
 - <u>Hash-based sharding</u>: Documents are distributed according to an MD5 hash \rightarrow uniform distribution
 - <u>Range-based sharding</u>: Documents with shard key values close to one another are likely to be co-located on the same shard → works well for range queries
 - <u>Location-based sharding</u>: Documents are partitioned w.r.t. a user-specified configuration that associates shard key ranges with specific shards and hardware

Document Stores:



MongoDB Consistency & Availabilty:

- default: strong consistency (but configurable)
- increased availability through replication
 - *replica sets* consist of one *primary* and multiple *secondary members*
 - MongoDB applies writes on the primary and then records the operations on the primary's oplog



Big Picture



- rows are identified by keys
- rows can have different numbers of columns (up to millions)
- order of rows depend on key values (locality is important!)
- multiple rows can be summarized to *families* (or *tablets*)
- multiple families can be summarized to a *key space*







- developed by Facebook, Apache project since 2009
- cluster Architecture:
 - P2P system (ordered as rings)
 - Each node plays the same role (decentralized)



- Each node accepts read/write operations
- user access through nodes via Cassandra Query Language (CQL)

Wide Column Stores:



Consistency

- tunable Data Consistency (chosable per operation)
- read repair: if stale data is read, Cassandra issues a read repair \rightarrow find most up-to-date data and update stale data
- generally: Eventually consistent
- main focus on availability!

Big Picture



Graph Databases:

- use graphs to store and represent relationships between entities
- composed of *nodes* and *edges*
- each node and each edge can contain properties (Property-Graphs)



Graph Databases:



Alice is a friend of Bob and vice versa. They both love the movie "Titanic".

Graph Databases:



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Graph Databases:



Alice is a friend of Bob and vice versa. They both love the movie "Titanic".



Graph Databases:



- master-slave replication (no partitioning!)
- consistency: eventual consistency (tunable to Immediate consistency)
- support of ACID Transactions
- cypher Query Language
- schema-optional

Big Picture

